

The difference between the weak and the strong core from the design point of view

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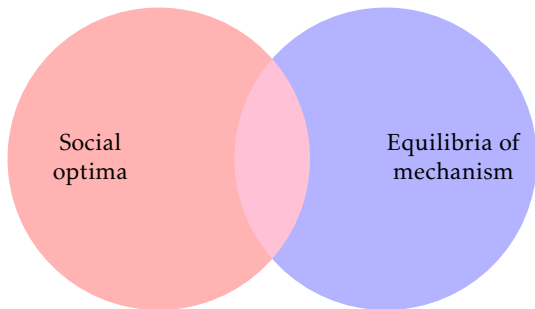
Mechanism design vs implementation

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- (Full) implementation: All equilibria are socially optimal **and** each social optimum can be realized through some equilibrium.



Implementation via mechanisms vs implementation via rights structures

Mechanisms	Rights structures
Strategies	Rights
Noncooperative equilibrium notions	Cooperative equilibrium notions

Mechanisms and rights structures, example

Primitives

$$N = \{1, 2\}, \mathcal{R} = \{R, R'\}$$

$$\phi(R) = \{x, w\}, \phi(R') = \{x\}$$

R_1	R_2	R'_1	R'_2
x	w	x	w
z	z	z	x
w	x	y	z
y	y	w	y

Table: Preferences

$1/2$	l	r
U	x	y
D	z	w

Table: Mechanism/game form

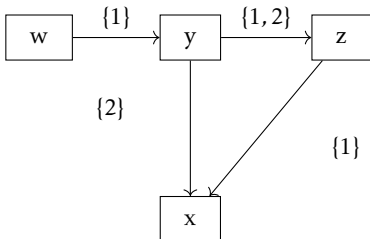


Figure: Rights structure

To cooperate or not?

R_1	R_2
xy	x
z	y
	z

Figure:
Preferences

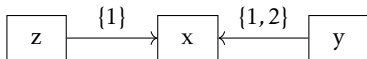


Figure: Rights structure

- Weak core: $C(\Gamma, R) = \{x, y\}$.
- Strong core: $SC(\Gamma, R) = \{x\}$.

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This paper

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- What SCRs can we implement in weak core?
- What SCRs can we implement in strong core?
- Remarks and applications.

Related literature

Noncooperative mechanisms
Maskin (1999) RES
Moore and Repullo (1990) JET
Jackson (1991) ECMA
Dutta and Sen (2012) GEB
De Clippel (2014) AER

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Rights structures

Koray and Yildiz (2018) JET
Korpela et al (2020) JET
Korpela et al (2021) GEB
Savva (2021) EL
Lombardi et al (2021)

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- Set of agents $N = \{1, 2, \dots, n\}$, set of all coalitions \mathcal{N} and set of all non-empty coalitions \mathcal{N}_0 .

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Interpretation: social planner wants to implement ϕ , but knows only \mathcal{R} .

Rights structures

Object of design for the social planner: means to implement SCR ϕ .

Key idea: allocation of rights in the society (constitution).

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Strong core: $s \in SC(\Gamma, R)$ if for all $t \in S$, there is no $K \in \gamma(s, t)$, such that for all $i \in K$ we have $h(t)R_i h(s)$ with strict preference for at least one $j \in K$.

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Strong core: $s \in SC(\Gamma, R)$ if for all $t \in S$, there is no $K \in \gamma(s, t)$, such that for all $i \in K$ we have $h(t)R_i h(s)$ with strict preference for at least one $j \in K$.

Implementation:

for any $R \in \mathcal{R}$, $\phi(R) = h(C(\Gamma, R))$ (or $= h(SC(\Gamma, R))$).

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Implementation in weak core

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Maskin-monotonicity

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There exists $Y \subseteq X$ with $\phi(\mathcal{R}) \subseteq Y$, such that for all R, R' and $x \in \phi(R)$:

if for all i , $[L_i(x, R) \subseteq L_i(x, R')] \cap Y$, then $x \in \phi(R')$.

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There exists $Y \subseteq X$ with $\phi(\mathcal{R}) \subseteq Y$, such that for all R and $x \in Y$:
if for all i , $Y \subseteq L_i(x, R)$, then $x \in \phi(R)$.

Theorem

The following are equivalent:

- (i) ϕ is implementable in weak core.
- (ii) ϕ satisfies Maskin-monotonicity and unanimity.
- (iii) ϕ is implementable in weak core by an individual-based rights structure.

Example

$$\phi(R) = \{x, w\}, \phi(R') = \{x\}$$

R_1	R_2	R'_1	R'_2
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Table: Preferences

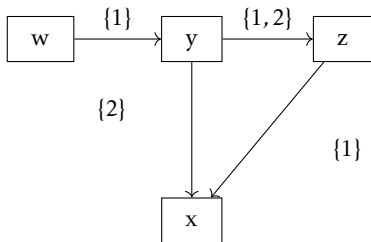


Figure: Rights structure

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- Whichever SCR is implementable by rights structures in weak core, is also implementable by an *individual-based* rights structure.
- There is no value in *coalitional rights* from the economic design point of view.

Implementation in strong core

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A few further definitions:

$$I_K(x, R) = \bigcap_{i \in K} \{y \in X \mid x I_i y \text{ for all } i \in K\}, \text{ and}$$
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Weak SC-monotonicity

There exists $Y \supseteq \phi(\mathcal{R})$ such that for all R, R' and $x \in \phi(R)$:

if for all K , $[I_K(x, R) \cup SL_K(x, R)] \cap Y \subseteq I_K(x, R') \cup SL_K(x, R')$, then $x \in \phi(R')$.

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Theorem

The following are equivalent:

- (i) *An SCR ϕ is implementable in strong core.*
- (ii) *An SCR ϕ satisfies weak SC-monotonicity and unanimity with respect to Y .*

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 - SC-monotonicity does not imply and is not implied by Maskin-monotonicity.
- Equivalence between individual-based and non individual-based rights structures breaks down.
 - Coalitions matter!
 - Restoring value of coalitional rights in mechanism design.

The rationing problem

- Lombardi et al (2022).
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- Set of allocations $Z \equiv \{x \in \mathbb{R}_+^n \mid \sum_{i \in N} x_i = M\}$.
- For all $i \in N$, $\succeq_i \in \mathcal{L}_i$ and $x, y \in Z$:
$$x R_i y \text{ if and only if } x_i \succeq_i y_i.$$

The no-envy correspondence

No-envy correspondence

$NE : \mathcal{R} \rightarrow 2^Z$, such that, for all $x \in Z$, we have $x \in NE(R)$ if and only if for all $\{i, j\} \subseteq N, x_i \succeq_i x_j$.

We are in general interested in the intersection of no-envy with strong (or weak) Pareto.

Proposition

Suppose that $NE \cap SPO$ is non-empty. Then, it is implementable in strong core.

Single-plateaued preferences

$N \equiv \{1, 2\}$, $\mathcal{R} = \{R, R'\}$, $M = 20$.

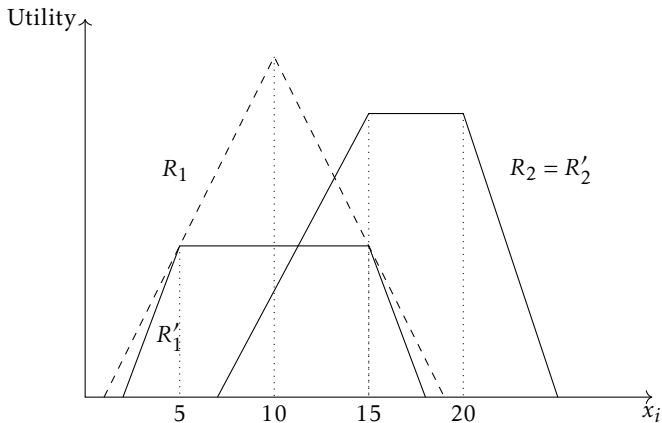


Figure: Single-plateaued preferences

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An open question

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Double implementation: for all $R \in \mathcal{R}$, $\phi(R) = h(C(\Gamma, R)) = h(SC(\Gamma, R))$.

Summary

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- Restoration of the value of coalitions in economic design.
- Importance of strong core as an equilibrium concept.

The end

Thanks! :-)

Rights structures example

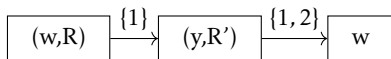


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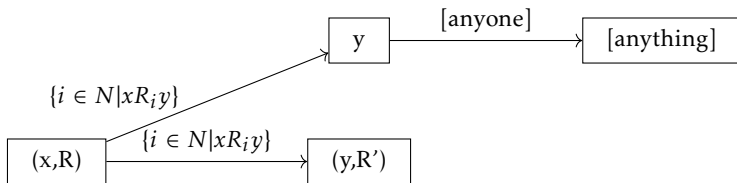


Figure: Canonical rights structure